Design and Analysis of Aging-Aware and Area Efficient Vedic Multiplier with Adaptive Hold Logic

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Abstract:
High speed, low power consumption is the key requirements to any VLSI design. The Area efficient multipliers play an important role. This paper presents an efficient implementation of a high speed, Vedic multiplier using aging aware technique and adaptive hold logic. This study presented the design and implementation of Vedic multipliers using XILINX. In this work, Modified Vedic multiplier is having least area. The Modified Vedic multiplier with adaptive hold logic and aging awareness make this efficient and also reliable. 32 bit signed multiplication and fractional multiplication is carried out and verified with around 10000 test patterns.

Keywords: Delay, Vedic Multiplier, Aging Aware, Adaptive hold.

I. Introduction
Multiplication is one of the most widely used arithmetic operations. Due to this a wide range of multiplier architectures are reported in the literature providing flexible choices for various applications. Among them the simplest is array multiplier [1] which is also the slowest. Some high performance multipliers are presented in [3-6]. Multiplication [4] is the most important arithmetic operation in signal processing applications. All the signal and data processing operations involve multiplication. As speed is always a constraint in the multiplication operation, increase in speed can be achieved by reducing the number of steps in the computation process. The speed of multiplier determines the efficiency of such a system. In any system design, the three main constraints which determine the performance of the system are speed, area and power requirement.

Vedic mathematics [2] was reconstructed from the ancient Indian scriptures (Vedas) by Swami Bharati Krishna Tirthaji Maharaja (1884-1960) after his eight years of research on Vedas. Vedic mathematics is mainly based on sixteen principles or word-formulae which are termed as sutras. This is a very interesting field and presents some effective algorithms which can be applied to various branches of engineering such as computing and digital signal processing. Integrating multiplication with Vedic Mathematics techniques would result in the saving of computational time.
Thus, integrating Vedic mathematics for the multiplier design will enhance the speed of multiplication operation.

II. Literature Survey

Traditional circuits use critical path delay as the overall circuit clock cycle in order to perform correctly. However, the probability that the critical paths are activated is low. In most cases, the path delay is shorter than the critical path. For these Noncritical paths, using the critical path delay as the overall cycle period will result in significant timing waste. Hence, the variable-latency design was proposed to reduce the timing waste of traditional circuits. The variable-latency design divides the circuit into two parts: 1) shorter paths and 2) longer paths. Shorter paths can execute correctly in one cycle, whereas longer paths need two cycles to execute. When shorter paths are activated frequently, the average latency of variable-latency designs is better than that of traditional designs. For example, several variable-latency adders were proposed using the speculation technique with error detection and recovery [8].

A short path activation function algorithm was proposed in [9] to improve the accuracy of the hold logic and to optimize the performance of the variable-latency circuit. An instruction scheduling algorithm was proposed in [10] to schedule the Operations on no uniform latency functional units and improve the performance of Very Long Instruction Word processors. In [11], a variable-latency pipelined multiplier architecture with a Booth algorithm was proposed.

In [12], process-variation tolerant architecture for arithmetic units was proposed, where the effect of process-variation is considered to increase the circuit yield. In addition, the critical paths are divided into two shorter paths that could be unequal and the clock cycle is set to the delay of the longer one. These research designs were able to reduce the timing waste of traditional circuits to improve performance, but they did not consider the aging effect and could not adjust themselves during the runtime. A variable-latency adder design that considers the aging effect was proposed in [13] and [14]. However, no variable-latency multiplier design that considers the aging effect and can adjust dynamically has been done.

III. Design of Vedic Multiplier

As discussed in the previous section, the critical path delay of the Multiplier depends on the computation time of adder and multiplier. In this section, we have discussed an efficient architecture for adder and multiplier using modified Ripple Carry adder.

The 32-bit multiplier is designed using four 16x16 Vedic multipliers which employ Urdhva Tiryagbhyam sutra and carry skip technique for partial product addition. The output of these Vedic multipliers is added by modifying the logic levels of ripple carry adder. Block diagram of the proposed 32x32 multiplier is illustrated in figure 1.

The 32-bit input sequence is divided into two 16-bit numbers and given as inputs to the 16-bit multiplier blocks (a[31:16] & b[31:16], a[15:0] & b[31:16], a[31:16] & b[15:0], a[15:0] & b[15:0]). The four multipliers used (in figure 1) are similar and give 32-bit intermediate products which are added using overlapping logic with the help of three modified parallel adders (ADDER-1, ADDER-2 and ADDER-3). The partial products obtained from the four multipliers are demarcated into four regions. The four LSB product bits P[15:0] are directly obtained from one of the multipliers. The output of the second and third multiplier block is added directly using ADDER-1 as the second and third region is overlapping. Then the higher order bit of first multiplier block.
is added to the overlapping sum using ADDER-2 which gives the product $P[31:16]$. Finally, MSB bits $P[63:32]$ are obtained by adding the fourth multiplier output to the carry from ADDER-1 (added at the fifth bit position) and higher order bits (acts as lower nibble of addend) of ADDER-3.

![Fig 1 Block diagram of 32x32 multiplier.](image)

**IV. Proposed Aging Aware Multiplier Architecture**

This section details the proposed aging-aware reliable multiplier design. It introduces the overall architecture and the functions of each component and also describes how to design AHL that adjusts the circuit when significant aging occurs.

![Fig 2 Proposed aging aware Vedic multiplier](image)

Fig 2 shows our proposed aging-aware multiplier architecture, which includes two m-bit inputs ($m$ is a positive number), one 2m-bit output, one Vedic multiplier, 2m 1-bit Razor flip-flops [16], and an AHL circuit. The inputs of the row-bypassing multiplier are the symbols in the parentheses. In the proposed architecture, Vedic multiplier can be examined by the number of zeros in either the multiplicand or multiplicator to predict whether the operation requires one cycle or two cycles to complete. When input patterns are random, the number of zeros or ones in the multiplicator and multiplicand follows a normal distribution. Therefore, using the number of zeros or ones as the judging criteria results in similar outcomes. Hence, the two aging-aware multipliers can be implemented using similar architecture, and the difference between the two bypassing multipliers lies in the input signals of the AHL. According to the bypassing selection in the column or row-bypassing multiplier, the input signal of the AHL in the architecture with the column-bypassing multiplier is the multiplicand, whereas that of the row-bypassing multiplier is the multiplicator. Razor flip-flops can be used to detect whether timing violations occur before the next input pattern arrives.

Fig. 3 shows the details of Razor flip-flops. A 1-bit Razor flip-flop contains a main flip-flop, shadow latch, XOR gate, and mux. The main flip-flop catches the execution result for the combination circuit using a normal clock signal, and the shadow latch catches the execution result using a delayed clock signal, which is slower than the normal clock signal. If the latched bit of the shadow latch is different from that of the main flip-flop, this means the path delay of the current operation exceeds the cycle period, and the main flip-flop catches an incorrect result.
If errors occur, the Razor flip-flop will set the error signal to 1 to notify the system to reexecute the operation and notify the AHL circuit that an error has occurred. We use Razor flip-flops to detect whether an operation that is considered to be a one-cycle pattern can really finish in a cycle. If not, the operation is reexecuted with two cycles.

Although the reexecution may seem costly, the overall cost is low because the reexecution frequency is low. More details for the Razor flip-flop can be found in [17]. The AHL circuit is the key component in the aging-aware variable-latency multiplier. Fig. 4 shows the details of the AHL circuit. The AHL circuit contains an aging indicator, two judging blocks, one mux, and one D flip-flop. The aging indicator indicates whether the circuit has suffered significant performance degradation due to the aging effect. The aging indicator is implemented in a simple counter that counts the number of errors over a certain amount of operations and is reset to zero at the end of those operations. If the cycle period is too short, the column-or row-bypassing multiplier is not able to complete these operations successfully, causing timing violations. These timing violations will be caught by the Razor flip-flops, which generate error signals. If errors happen frequently and exceed a predefined threshold, it means the circuit has suffered significant timing degradation due to the aging effect, and the aging indicator will output signal 1; otherwise, it will output 0 to indicate the aging effect is still not significant, and no actions are needed.

**Results**

**RTL Schematic**

![RTL Schematic of Aging Aware Vedic Multiplier](https://edupediapublications.org/journals/index.php/IJR/)

**B) Output Plot**

![Output Plot](https://edupediapublications.org/journals/index.php/IJR/)
This paper proposed an aging-aware multiplier design with the AHL. The multiplier is able to adjust the AHL to mitigate performance degradation due to increased delay. The experimental results show that our proposed architecture with 32×32 Vedic multiplier is Delay efficient. The Proposed aging-aware Vedic multiplier with Adaptive hold logic Architecture was synthesized and simulated in Xilinx and Modelsim software respectively.

**Conclusion**

This paper proposed an aging-aware multiplier design with the AHL. The multiplier is able to adjust the AHL to mitigate performance degradation due to increased delay. The experimental results show that our proposed architecture with 32×32 Vedic multiplier is Delay efficient. The Proposed aging-aware Vedic multiplier with Adaptive hold logic Architecture was synthesized and simulated in Xilinx and Modelsim software respectively.

**References**


**Table I** Comparison of Aging-Aware Multiplier with Row bypass, Column bypass and Vedic Techniques

<table>
<thead>
<tr>
<th>Architecture</th>
<th>Memory (Kb)</th>
<th>Delay (ns)</th>
</tr>
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<tbody>
<tr>
<td>Aging-Aware Row Multiplier</td>
<td>270412</td>
<td>71.591</td>
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<tr>
<td>Aging-Aware Column Multiplier</td>
<td>230836</td>
<td>69.398</td>
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<tr>
<td>Aging-Aware Vedic Multiplier</td>
<td>184504</td>
<td>32.146</td>
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